

## Test B – Solutions

Name: \_\_\_\_\_ Number: \_\_\_\_\_  
 (as on attendance sheets)

## Calculators may be used for simple arithmetic operations only!

1. (12 pts.)  $A$  and  $B$  are sets with  $|B| = 5$ . There are exactly  $32768 = 2^{15}$  relations from  $A$  to  $B$ .

(a) What is  $|A \times B|$  (the cardinality of the Cartesian product)?

$$15 = \log_2(2^{15}).$$

(b) What is  $|A|$ ?

$$\frac{|A \times B|}{|B|} = 3.$$

2. (16 pts.) One of these formulas is valid, the other is not. Prove the correct one and provide a counterexample for the other. *Suggestion:* For the counterexample, let  $p(x, y)$  be an inequality in the universe of real numbers.

(a)  $\forall x \exists y p(x, y) \rightarrow \exists y \forall x p(x, y)$

**Not valid.** Let  $p(x, y)$  be  $x < y$ . Then the hypothesis is true, because, given  $x$ ,  $y$  could be  $x + 1$ . However, the conclusion is false, because, whatever  $y$  is, one can find an  $x$  that is larger, say  $x = y + 1$ .

(b)  $\exists y \forall x p(x, y) \rightarrow \forall x \exists y p(x, y)$

**Valid.** Here is a formal deduction:

- (1)  $\exists y \forall x p(x, y)$  (hypothesis)
- (2) Let  $c$  be such a  $y$ :  $\forall x p(x, c)$  (existential specification)
- (3) Let  $d$  be arbitrary:  $p(d, c)$  (universal specification)
- (4)  $\exists y p(d, y)$  (existential generalization)
- (5) Since  $d$  was arbitrary,  $\forall x \exists y p(x, y)$  (universal generalization)
- (6) Therefore,  $\exists y \forall x p(x, y) \rightarrow \forall x \exists y p(x, y)$  (because we've shown that (1) implies (5))

**Remark:** Let's see why this type of argument fails in case (a):

- (1)  $\forall x \exists y p(x, y)$  (hypothesis)
- (2) Let  $d$  be arbitrary:  $\exists y p(d, y)$  (universal specification)
- (3) Let  $c$  be such a  $y$ :  $p(d, c)$  (existential specification)
- (4)  $\forall x p(x, c)$  (universal generalization??) **This is wrong, because in (3)  $c$  depends on  $x$ !** (See W. V. Quine, *Methods of Logic*, for instructions on how to use "flagged variables" to avoid this kind of error.)

3. (21 pts.) Recall that a pair of dice consists of two cubes, the sides of each of which are labeled by the numbers 1 through 6. When the dice are thrown, what is the probability of obtaining each of these results? (Leave the answer as a fraction in lowest terms, not a decimal.)

(a) a pair (i.e, the same number on each die)

Note first that there are  $6^2 = 36$  total possibilities, since the 6 possibilities for each die are independent of those for the other die. Of these, 6 outcomes are pairs, so the probability of a pair is  $\frac{6}{36} = \frac{1}{6}$ .

(b) a total of 4

There are 3 ways to get a 4:  $1 + 3$ ,  $2 + 2$ , or  $3 + 1$ . So the probability is  $\frac{3}{36} = \frac{1}{12}$ .

(c) either a total of 4, or a pair

All of the cases found in (a) and (b) count, but the case  $2 + 2$  must not be counted twice. So the total number of good outcomes is  $6 + 3 - 1 = 8$ , for a probability of  $\frac{8}{36} = \frac{2}{9}$ . This is an instance of the principle

$$|A \cup B| = |A| + |B| - |A \cap B|.$$

4. (16 pts.) Establish the validity of the argument

$$\{(p \rightarrow q) \wedge [(q \wedge r) \rightarrow s] \wedge r\} \Rightarrow (p \rightarrow s).$$

(Write a formal deduction, not a truth table.)

- (1)  $\{(p \rightarrow q) \wedge [(q \wedge r) \rightarrow s] \wedge r\}$  (hypothesis)
- (2)  $p$  (hypothesis of the conclusion)
- (3)  $(p \rightarrow q)$  (part of (1))
- (4)  $q$  (modus ponens, (2) and (3))
- (5)  $r$  (part of (1))
- (6)  $q \wedge r$  (from (4) and (5))
- (7)  $(q \wedge r) \rightarrow s$  (part of (1))
- (8)  $s$  (modus ponens, (6) and (7))
- (9)  $p \rightarrow s$  (because we have shown that (2) implies (8))

This shows that the hypothesis (1) logically implies the conclusion (9).

5. (15 pts.)

(a) Prove that  $A \cup (A \cap B) = A \cap (A \cup B)$ , or give a counterexample.

Method 1: Apply the distributive law for sets:

$$A \cup (A \cap B) = (A \cup A) \cap (A \cup B) = A \cap (A \cup B).$$

(Or apply the other distributive law from the other direction.)

Method 2: Appeal to the absorption law (part (b)) to show that both sides are equal to  $A$ .

Method 3: An “element proof”:

$$x \in A \cup (A \cap B) \Rightarrow (x \in A) \vee (x \in A \cap B).$$

But  $x \in A \cap B \rightarrow x \in A$ . So in either case of “ $\vee$ ” we have  $x \in A$ . And if  $x \in A$ , then also  $x \in A \cup B$ . And therefore  $x \in A \cap (A \cup B)$ . So far we have shown that

$$A \cup (A \cap B) \subseteq A \cap (A \cup B).$$

Conversely, assume that  $x \in A \cap (A \cup B)$ . Then  $x \in A$  and  $x \in A \cup B$ . But if  $x \in A$ , then certainly  $x \in A \cup (A \cap B)$ , and we have now shown that

$$A \cap (A \cup B) \subseteq A \cup (A \cap B).$$

So the two sets are equal.

Method 4: Start from the logical absorption law,  $p \wedge (p \vee q) \longleftrightarrow p \vee (p \wedge q)$ , or the logical distributive law,  $p \wedge (q \vee r) \longleftrightarrow (p \wedge q) \vee (p \wedge r)$ ; let  $p$  be  $x \in A$ , etc., and apply the definitions of the set-theory operations.

(b) Can you simplify either side of the equation to something else? Yes, both sides simplify to  $A$ . This is called the “absorption law” for sets.

6. (20 pts.) Prove by induction that (for  $n = 0, 1, \dots$ )

$$\frac{d^n}{dx^n}[f(x)g(x)] = \sum_{j=0}^n \binom{n}{j} f^{(j)}(x)g^{(n-j)}(x).$$

(Here  $f^{(j)}$  is the  $j$ th derivative of  $f$ , etc.) *Hints:* Use  $\binom{n-1}{j-1} + \binom{n-1}{j} = \binom{n}{j}$ . Start by showing that the case  $n = 1$  is a well known fact of calculus.

The case  $n = 1$  is the product rule,  $\frac{d}{dx}[f(x)g(x)] = f(x)g'(x) + f'(x)g(x)$ . The case  $n = 0$  is trivial:  $f(x)g(x) = f(x)g(x)$ . We need to see what happens when we apply the product rule repeatedly. Let's write out the formula for the case  $n = 1$ :

$$\frac{d^{n-1}}{dx^{n-1}}(fg) = \sum_{j=0}^{n-1} \binom{n-1}{j} f^{(j)}g^{(n-1-j)}.$$

Then

$$\begin{aligned} \frac{d^n}{dx^n}(fg) &= \frac{d}{dx}(fg)^{(n-1)} \\ &= \sum_{j=0}^{n-1} \binom{n-1}{j} \left[ f^{(j+1)}g^{(n-1-j)} + f^{(j)}g^{(n-j)} \right] \\ &= \sum_{j=1}^n \binom{n-1}{j-1} f^{(j)}g^{(n-j)} + \sum_{j=0}^{n-1} \binom{n-1}{j} f^{(j)}g^{(n-j)} \\ &= \binom{n-1}{0} f^{(0)}g^{(n)} + \binom{n-1}{n-1} f^{(n)}g^{(0)} + \sum_{j=1}^{n-1} \left[ \binom{n-1}{j-1} + \binom{n-1}{j} \right] f^{(j)}g^{(n-j)} \\ &= f^{(0)}g^{(n)} + f^{(n)}g^{(0)} + \sum_{j=1}^{n-1} \binom{n}{j} f^{(j)}g^{(n-j)} \\ &= \sum_{j=0}^n \binom{n}{j} f^{(j)}g^{(n-j)}. \end{aligned}$$

**Remark:** The hint is the identity on which Pascal's triangle is based. So the theorem says that repeated differentiation involves exactly the same combinatorics as the binomial theorem.